

# More on Event-B: Relations

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# Ordered Pairs and Cartesian Products

An **ordered pair** is an element consisting of two parts:  
a **first** part and a **second** part.

An ordered pair with first part  $x$  and second part  $y$  is written:  $x \mapsto y$

The **Cartesian product** of two sets is the **set of pairs** whose first part is in  $S$  and second part is in  $T$ .

The Cartesian product of  $S$  with  $T$  is written:  $S \times T$

# Cartesian Products: Definition and Examples

Defining Cartesian product:

Predicate	Definition
$x \mapsto y \in S \times T$	$x \in S \wedge y \in T$

Examples:

$$\{a, b, c\} \times \{1, 2\} = \{ a \mapsto 1, a \mapsto 2, b \mapsto 1, \\ b \mapsto 2, c \mapsto 1, c \mapsto 2, \}$$

$$\{a, b, c\} \times \{\} = ?$$

$$\{ \{a\}, \{a, b\} \} \times \{1, 2\} = ?$$

# Cartesian Product is a Type Constructor

$S \times T$  is a new type constructed from types  $S$  and  $T$ .

Cartesian product is the type constructor for ordered pairs.

Given  $x \in S$ ,  $y \in T$ , we have

$$x \mapsto y \in S \times T$$

$$4 \mapsto 7 \in ?$$

$$\{5, 6, 3\} \mapsto 4 \in ?$$

$$\{4 \mapsto 8, 3 \mapsto 0, 2 \mapsto 9\} \in ?$$

# Sets of Order Pairs

A database can be modelled as a **set of ordered pairs**:

$$\begin{aligned} \textit{directory} = \{ & \textit{mary} \mapsto 287573, \\ & \textit{mary} \mapsto 398620, \\ & \textit{john} \mapsto 829483, \\ & \textit{jim} \mapsto 398620 \} \end{aligned}$$

*directory* has type

$$\textit{directory} \in \mathbb{P}(\textit{Person} \times \textit{PhoneNum})$$

# Relations

A **relation** is a set of ordered pairs.

A relation is a common modelling structure so Event-B has a special notation for it:

$$\boxed{T \leftrightarrow S} = \mathbb{P}(T \times S)$$

So we can write:

$$directory \in Person \leftrightarrow PhoneNum$$

Do not confuse the arrow symbols:

$\leftrightarrow$  combines **two sets** to form a **set**.

$\mapsto$  combines **two elements** to form an **ordered pair**.

# Domain and Range

- ▶ The **domain** of a relation  $R$  is the set of first parts of all the pairs in  $R$ , written  $\boxed{dom(R)}$
- ▶ The **range** of a relation  $R$  is the set of second parts of all the pairs in  $R$ , written  $\boxed{ran(R)}$

Examples:

$$dom(directory) = \{mary, john, jim\}$$

$$ran(directory) = \{287573, 398620, 829483\}$$

## Domain and Range Definitions

Predicate	Definition
$x \in \text{dom}(R)$	$\exists y \cdot x \mapsto y \in R$
$y \in \text{ran}(R)$	$\exists x \cdot x \mapsto y \in R$



# Telephone Directory Model

- ▶ Phone directory relates people to their phone numbers.
- ▶ Each person can have zero or more numbers.
- ▶ People can share numbers.

```
context PhoneContext  
sets Person PhoneNum  
end
```

```
machine PhoneBook  
variables dir  
invariants  $dir \in Person \leftrightarrow PhoneNum$ 
```

```
initialisation  $dir := \{\}$ 
```

## Extending the Directory

Add an entry to the directory:

```
AddEntry  $\hat{=}$  any  $p, n$  where  
     $p \in Person$   
     $n \in PhoneNum$   
then  
     $dir := dir \cup \{p \mapsto n\}$   
end
```

## Relational Image

Assume  $R \in S \leftrightarrow T$  and  $A \subseteq S$

The **relational image** of set  $A$  under relation  $R$  is written

$R[A]$

Predicate	Definition
$y \in R[A]$	$\exists x \cdot x \in A \wedge x \mapsto y \in R$

Example:

$$\begin{aligned} \text{directory} = \{ & \text{mary} \mapsto 287573, \\ & \text{mary} \mapsto 398620, \\ & \text{john} \mapsto 829483, \\ & \text{jim} \mapsto 398620 \} \end{aligned}$$
$$\text{directory}[\{\text{mary}\}] = \{287573, 398620\}$$

# Modelling Queries using Relational Image

Determine all the numbers associated with a person in the directory:

```
GetNumbers  $\hat{=}$  any p, result where  
    p  $\in$  Person  
    result = dir[ {p} ]  
end
```

Determine all the numbers associated with a set of people:

```
GetMultiNumbers  $\hat{=}$  any ps, result where  
    ps  $\subseteq$  Person  
    result = dir[ ps ]  
end
```

## Relational Inverse

Given  $R \in S \leftrightarrow T$ , the **relational inverse** of  $R$  is written  $R^{-1}$

Predicate	Definition
$y \mapsto x \in R^{-1}$	$x \mapsto y \in R$

Example:

$$\begin{aligned} \text{directory}^{-1} = \{ & 287573 \mapsto \text{mary}, \\ & 398620 \mapsto \text{mary}, \\ & 829483 \mapsto \text{john}, \\ & 398620 \mapsto \text{jim} \} \end{aligned}$$

$$\text{directory}^{-1}[\{398620\}] = \{\text{mary}, \text{jim}\}$$

## Inverse Queries

Return all the people associated with a number in the directory:

```
GetNames  $\hat{=}$  any n, result where  
    n  $\in$  PhoneNum  
    result = dir-1[ {n} ]  
end
```

Return all the people associated with a set of numbers:

```
GetMultiNames  $\hat{=}$  any ns, result where  
    ns  $\subseteq$  PhoneNum  
    result = dir-1[ ns ]  
end
```

# Domain Restriction

Given  $R \in S \leftrightarrow T$  and  $A \subseteq S$ ,  
the **domain restriction** of  $R$  by  $A$  is written  $A \triangleleft R$

Restrict relation  $R$  so that it only contains pairs whose first part is in the set  $A$ .

Example:

$$\text{directory} = \{ \text{mary} \mapsto 287573, \text{mary} \mapsto 398620, \\ \text{john} \mapsto 829483, \text{jim} \mapsto 398620 \}$$

$$\{ \text{john}, \text{jim}, \text{jane} \} \triangleleft \text{directory} = \{ \text{john} \mapsto 829483, \\ \text{jim} \mapsto 398620 \}$$

# Domain Subtraction

Given  $R \in S \leftrightarrow T$  and  $A \subseteq S$ ,  
the **domain subtraction** of  $R$  by  $A$  is written  $A \triangleleft R$

Remove those pairs from  $R$  whose first part is in  $A$ .

Example:

$$\{john, jim, jane\} \triangleleft directory = \{ mary \mapsto 287573, \\ mary \mapsto 398620 \}$$



# Domain and Range, Restriction and Subtraction

Assume  $R \in S \leftrightarrow T$  and  $A \subseteq S$  and  $B \subseteq T$

Predicate	Definition	
$x \mapsto y \in A \triangleleft R$	$x \mapsto y \in R \wedge x \in A$	domain restriction
$x \mapsto y \in A \triangleleft R$	$x \mapsto y \in R \wedge x \notin A$	domain subtraction
$x \mapsto y \in R \triangleright B$	$x \mapsto y \in R \wedge y \in B$	range restriction
$x \mapsto y \in R \triangleright B$	$x \mapsto y \in R \wedge y \notin B$	range subtraction

## Removing Entries from the Directory

Remove all the entries associated with a person in the directory:

```
RemovePerson  $\hat{=}$  any p where  
    p  $\in$  Person  
then  
    dir := {p}  $\triangleleft$  dir  
end
```

Remove all the entries associated with a number in the directory:

```
RemoveNumber  $\hat{=}$  any n where  
    n  $\in$  PhoneNum  
then  
    dir := dir  $\triangleright$  {n}  
end
```

# Relational Composition

Given  $Q \in S \leftrightarrow T$  and  $R \in T \leftrightarrow U$ ,  
the **relational composition** of  $Q$  and  $R$  is written  $Q ; R$

We have that  $Q ; R \in S \leftrightarrow U$

Predicate	Definition
$x \mapsto z \in (Q ; R)$	$\exists y \cdot x \mapsto y \in Q \wedge y \mapsto z \in R$

Example:

$$M = \{ a \mapsto l, b \mapsto m, c \mapsto n \}$$

$$N = \{ l \mapsto 4, n \mapsto 6, p \mapsto 8 \}$$

$$M ; N = ?$$

## Composition and Image

Given  $Q \in S \leftrightarrow T$  and  $R \in T \leftrightarrow U$  and  $A \subseteq S$

$$(Q ; R)[A] = R[Q[A]]$$

Example:

$$M = \{ a \mapsto l, b \mapsto m, c \mapsto n \}$$

$$N = \{ l \mapsto 4, n \mapsto 6, p \mapsto 8 \}$$

$$(M ; N) [ \{a, b\} ] = ?$$

$$N[ M[ \{a, b\} ] ] = ?$$

## Extend directory with friends

**variables**  $dir, friend$

**invariants**

$friend \in Person \leftrightarrow Person$

$dir \in Person \leftrightarrow PhoneNum$

Return the telephone numbers of all friends of  $p$ :

$GetFriendNumbers \hat{=}$

**any**  $p, result$  **where**

$p \in Person$

$result = (friend; dir)[ \{p\} ]$

**then**

$skip$

**end**

# Recap

- ▶ Cartesian product is the type constructor for pairs of elements.
- ▶ A relation is a set of pairs.
- ▶ Range of a relation, domain of a relation.
- ▶ Relational image, relational inverse.
- ▶ Restriction and subtraction.
- ▶ Relational composition.