





Efficient Budget Allocation with Accuracy Guarantees for Crowdsourcing Classification Tasks

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1. Motivation Crowdsourcing systems:

Churth

₿C

Use mass of human intelligence to solve hard AI problems



amazon mechanical turk Artificial Artificial Intelligence

Single response is *unreliable redundant task allocation*

How many responses/task is *sufficient*?

- users are paid for their responses
- restricted budget limit
- ElanceOnline

amazon mechanical turk^{**}

- budget has to be assigned to tasks at the beginning (online approaches are not suitable here)
- different tasks may provide different payments

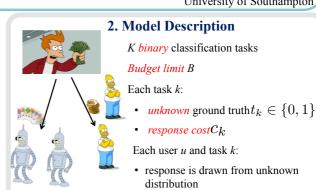
4. Performance Analysis

THEOREM 1 (MAIN RESULT 1). For any $B \ge \sum_{j=1}^{K} c_j$, the expected estimation error of CrowdBudget is bounded as follows:

$$\mathbb{E}\left[\Delta\left(B\right)\right] \le K e^{\frac{-2Bd_{\min}^2}{c_{\max}^2 \sum_{j=1}^{K} \frac{1}{c_j}}}.$$

THEOREM 2 (MAIN RESULT 2). For any $B \ge \sum_{j=1}^{K} c_j$ and $0 < \beta < 1$, the following bound holds for the expected es-timation error of CrowdBudget with at least $(1 - \beta)^K$ probability:

$$\mathbb{E}\left[\Delta\left(B\right)\right] \le \max\left\{0, \frac{K}{2} - d_{\min}\sqrt{\frac{2B}{-\ln\frac{\beta}{2}}\sum_{j=1}^{K}\frac{1}{c_j}}\right\}$$



We assume that the expected response is *closer* to the ground truth

$$\mathbb{E}_{X_k}\left[\mu_{u,k}\right] = \mu_k \quad \forall k : |\mu_k - t_k| < \frac{1}{2}$$

Let $\hat{r}(n_k)$ denote the *estimate* of the truth

• derived from the received responses by using a *fusion method*

MV-efficient method: provides estimation at least as good as majority voting

Estimation error $\Delta(n_k) = |\hat{r}_k(n_k) - t_k|$ Total exp. error $\mathbb{E}[\Delta(B)] = \sum_{k=1}^{K} \mathbb{E}[\Delta(n_k)] = \sum_{k=1}^{K} P\{\hat{r}_k(n_k) \neq t_k\}$

Objective

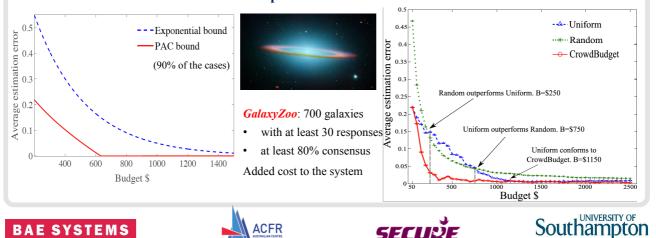
Minimise the total expected error, s.t. budget limit B Provide theoretical performance guarantees

3. Algorithm

Step 1:
$$n_k = \left\lfloor \frac{B}{c_k^2 \sum_{j=1}^K \frac{1}{c_j}} \right\rfloor$$

Step 2: while $B^r \ge 0$ and $k \le K$ do $n_k = n_k + 1, B^r = B^r - c_k, k = k + 1$

Step 3: Use an MV-efficient method to estimate responses



5. Empirical Evaluation